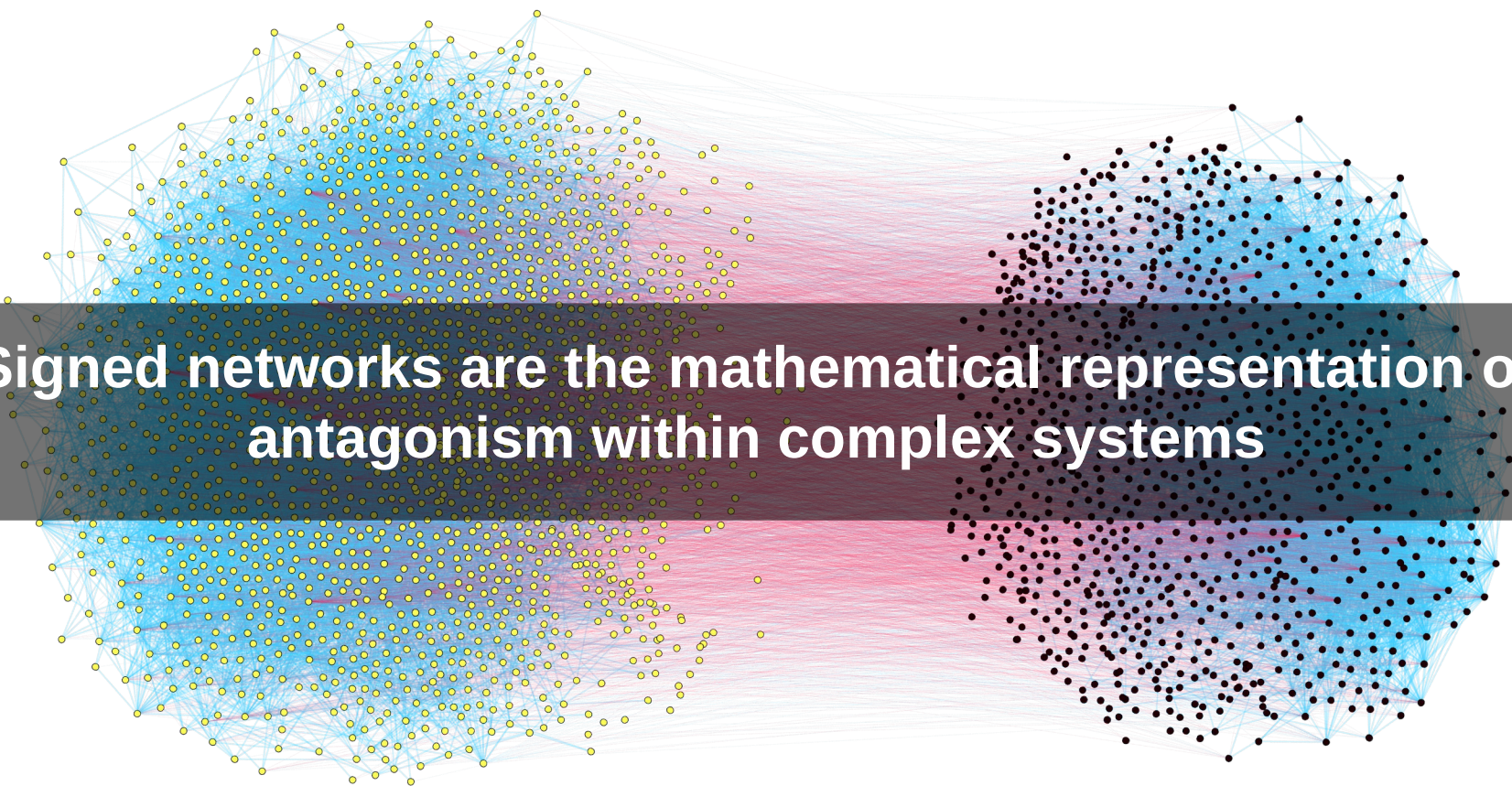


Friends, Foes, and Factions: The math behind signed networks

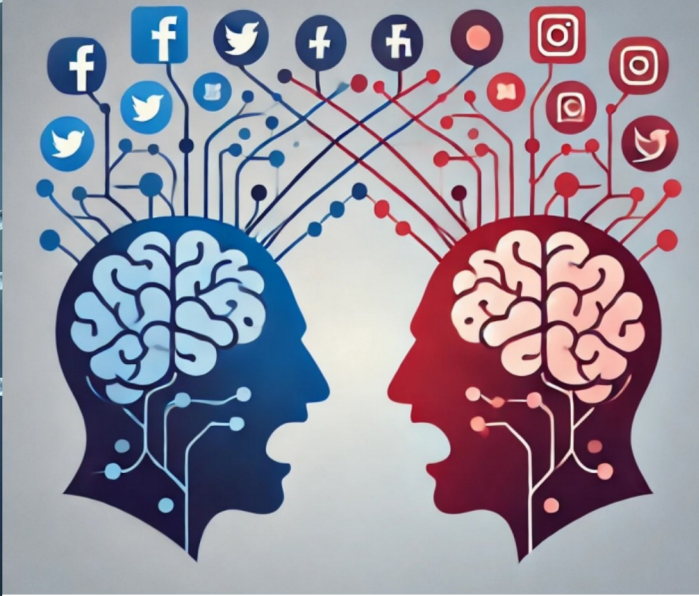
Fernando Díaz Díaz

**Pisan Young Seminars in Applied and NUmberical
Mathematics (PYSANUM)**





Signed networks are the mathematical representation of antagonism within complex systems





LETTER

nature

Stability criteria for complex ecosystems

Stefano Allesina^{1,2} & Si Tang¹



LETTER

nature

Stability criteria for complex ecosystems

Stefano Allesina^{1,2} & Si Tang¹

BMC Systems Biology

BioMed Central

Research article

Open Access

The regulatory network of *E. coli* metabolism as a Boolean dynamical system exhibits both homeostasis and flexibility of response

Areejit Samal¹ and Sanjay Jain^{*1,2,3}



The role of negative edges



ARTICLE

<https://doi.org/10.1038/s41467-019-10548-8>

OPEN

Structural balance emerges and explains performance in risky decision-making

Omid Askarisichani¹, Jacqueline Ng Lane², Francesco Bullo^{3,4}, Noah E. Friedkin^{3,5}, Ambuj K. Singh¹ & Brian Uzzi^{6,7}

The role of negative edges



ARTICLE

<https://doi.org/10.1038/s41467-019-10548-8>

OPEN

Structural balance emerges and explains performance in risky decision-making

Omid Askarisichani¹, Jacqueline Ng Lane², Francesco Bullo^{3,4}, Noah E. Friedkin^{3,5}, Ambuj K. Singh¹ & Brian Uzzi^{6,7}

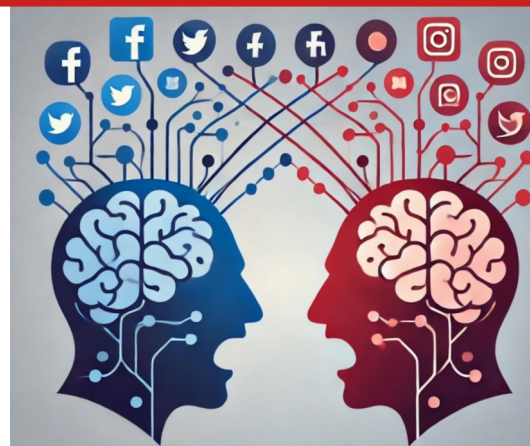
PNAS

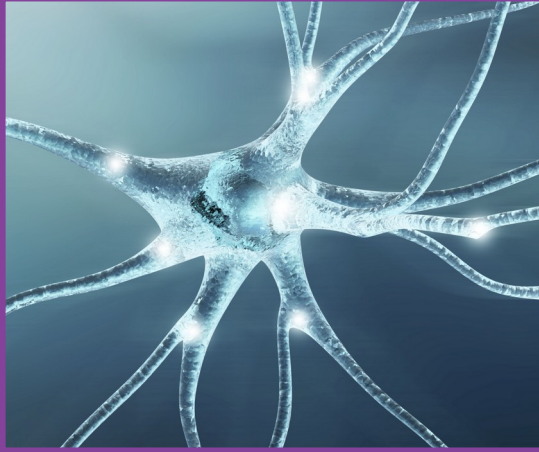
RESEARCH ARTICLE

| SOCIAL SCIENCES

Triadic influence as a proxy for compatibility in social relationships

Miguel Ruiz-García  ^{a,b,c,2,1}, Juan Ozaita  ^{c,1}, María Pereda  ^{b,d}, Antonio Alfonso^e, Pablo Brañas-Garza  ^e, José A. Cuesta  ^{b,c,f}, and Angel Sánchez  ^{b,c,f}





The human brain is intrinsically organized into dynamic, anticorrelated functional networks

Michael D. Fox^{*†}, Abraham Z. Snyder^{*‡}, Justin L. Vincent^{*}, Maurizio Corbetta[‡], David C. Van Essen⁵, and Marcus E. Raichle^{*‡§¶}



The human brain is intrinsically organized into dynamic, anticorrelated functional networks

Michael D. Fox^{*†}, Abraham Z. Snyder^{*‡}, Justin L. Vincent^{*}, Maurizio Corbetta[‡], David C. Van Essen[§], and Marcus E. Raichle^{*‡§¶}

Power and Interdependence

ROBERT O. KEOHANE AND JOSEPH S. NYE, JR



Structural balance & signed metrics

Context X

Context Y

Mappings to unsigned networks

G

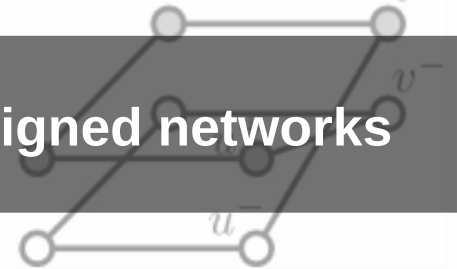
v



\mathcal{G}

v^+

v

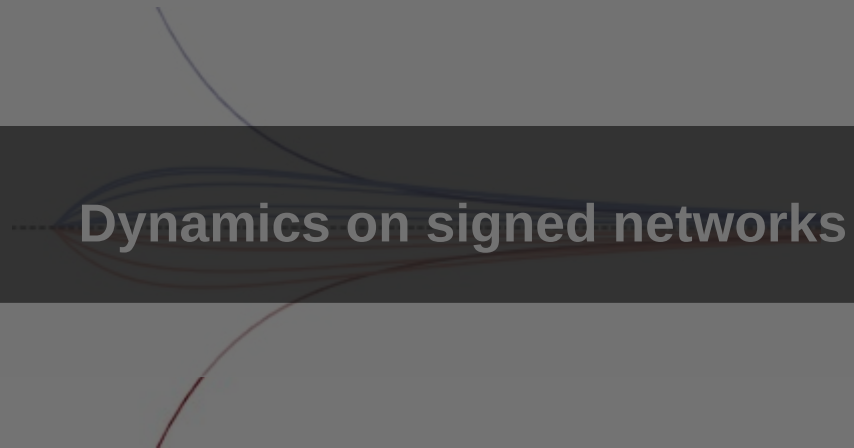
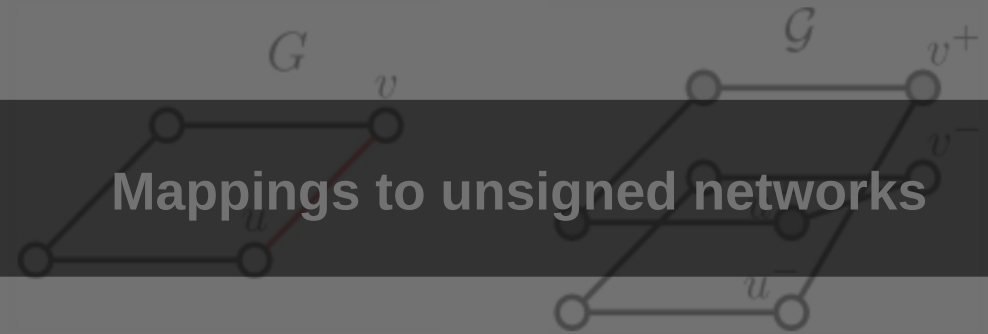


Dynamics on signed networks

Detection of communities and factions

(a)

(b)



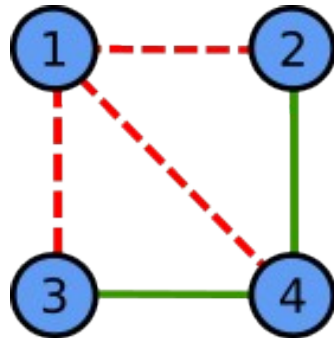
Mathematical representation of a signed network

Definition 1: a **signed network** is a triple $G = (V, E, \sigma)$, where V is the set of nodes, E is the set of edges, and $\sigma : E \rightarrow \{+1, -1\}$ is a function that assigns a sign (+ or -) to each edge.

Mathematical representation of a signed network

Definition 1: a **signed network** is a triple $G = (V, E, \sigma)$, where V is the set of nodes, E is the set of edges, and $\sigma : E \rightarrow \{+1, -1\}$ is a function that assigns a sign (+ or -) to each edge.

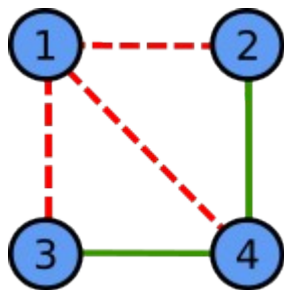
Definition 2: the **adjacency matrix** A of a signed network is a square matrix with elements $A_{ij} = \sigma(e)$ when $e=(i,j)$, and zero otherwise.



$$A = \begin{pmatrix} 0 & -1 & -1 & -1 \\ -1 & 0 & 0 & 1 \\ -1 & 0 & 0 & 1 \\ -1 & 1 & 1 & 0 \end{pmatrix}$$

The walk lemma for signed networks

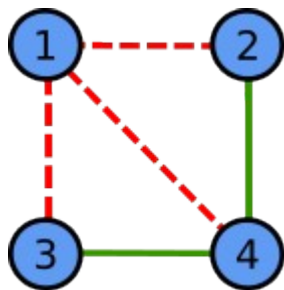
Definition 3: a **walk** is an ordered sequence of (not necessarily different) edges, where consecutive edges are incident to the same node. The **sign** of a walk is the product of the signs of its edges.



The walk lemma for signed networks

Definition 3: a **walk** is an ordered sequence of (not necessarily different) edges, where consecutive edges are incident to the same node. The **sign** of a walk is the product of the signs of its edges.

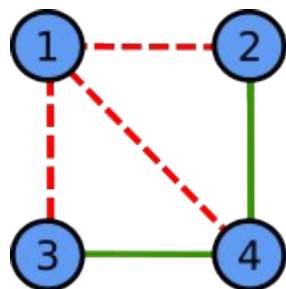
Lemma 4 (walk lemma): the (i,j) th element of the k -th power of the adjacency matrix, $(A^k)_{ij}$, counts the difference between the number of positive walks and the number of negative walks of length k between nodes i and j .



The walk lemma for signed networks

Definition 3: a **walk** is an ordered sequence of (not necessarily different) edges, where consecutive edges are incident to the same node. The **sign** of a walk is the product of the signs of its edges.

Lemma 4 (walk lemma): the (i,j) th element of the k -th power of the adjacency matrix, $(A^k)_{ij}$, counts the difference between the number of positive walks and the number of negative walks of length k between nodes i and j .



$$A^3 = \begin{pmatrix} 4 & -5 & -5 & -5 \\ -5 & 2 & 2 & 5 \\ -5 & 2 & 2 & 5 \\ -5 & 5 & 5 & 4 \end{pmatrix}$$

$$W_1 = \{e_{12}, e_{24}, e_{43}\}$$

$$W_2 = \{e_{12}, e_{21}, e_{13}\}$$

$$W_3 = \{e_{14}, e_{41}, e_{13}\}$$

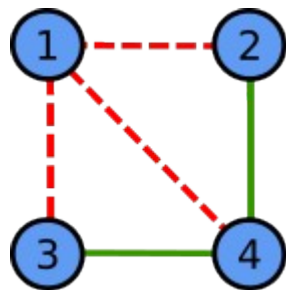
$$W_4 = \{e_{13}, e_{34}, e_{43}\}$$

$$W_5 = \{e_{13}, e_{31}, e_{13}\}$$

The walk lemma for signed networks

Definition 3: a **walk** is an ordered sequence of (not necessarily different) edges, where consecutive edges are incident to the same node. The **sign** of a walk is the product of the signs of its edges.

Lemma 4 (walk lemma): the (i,j) th element of the k -th power of the adjacency matrix, $(A^k)_{ij}$, counts the difference between the number of positive walks and the number of negative walks of length k between nodes i and j .



$$A^3 = \begin{pmatrix} 4 & -5 & -5 & -5 \\ -5 & 2 & 2 & 5 \\ -5 & 2 & 2 & 5 \\ -5 & 5 & 5 & 4 \end{pmatrix}$$

$$W_1 = \{e_{12}, e_{24}, e_{43}\}$$

$$W_2 = \{e_{12}, e_{21}, e_{13}\}$$

$$W_3 = \{e_{14}, e_{41}, e_{13}\}$$

$$W_4 = \{e_{13}, e_{34}, e_{43}\}$$

$$W_5 = \{e_{13}, e_{31}, e_{13}\}$$



COMBINATORICS



ALGEBRA



Harary's theory

Definition 5: A cycle or closed walk is **balanced** or positive if it contains an even number of negative edges.

-

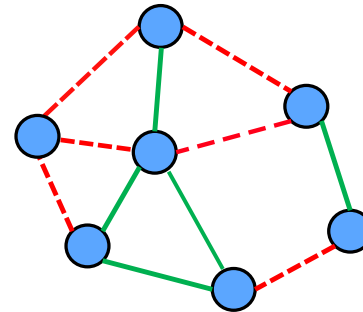
- Harary, F. (1953). "On the notion of balance of a signed graph". *Michigan Mathematical Journal*, 2(2), 143-146.
- Cartwright, D., & Harary, F. (1956). "Structural balance: a generalization of Heider's theory". *Psychological review*, 63(5), 277.

Harary's theory

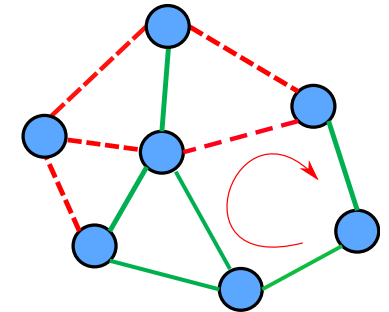
Definition 5: A cycle or closed walk is **balanced** or **positive** if it contains an even number of negative edges.

Definition 6: A signed network is **balanced** if every cycle within it is balanced. Otherwise, the network is **unbalanced**.

Balanced graph



Unbalanced graph



-

- Harary, F. (1953). "On the notion of balance of a signed graph". *Michigan Mathematical Journal*, 2(2), 143-146.
- Cartwright, D., & Harary, F. (1956). "Structural balance: a generalization of Heider's theory". *Psychological review*, 63(5), 277.

Balance theorems

Theorem 7 (Harary): a graph is balanced iff the node set can be split into two **balanced factions**, such that:

- Links within each balanced faction are positive, and
- Links between different balanced factions are negative.

- Harary, F. (1953). "On the notion of balance of a signed graph". *Michigan Mathematical Journal*, 2(2), 143-146.
- Acharya (1980). "Spectral criterion for cycle balance in networks". *Journal of Graph Theory*, 4(1), 1-11.
- Zaslavsky (1982). "Signed graphs". *Discrete Applied Mathematics* 4, 47-74.

Balance theorems

Theorem 7 (Harary): a graph is balanced iff the node set can be split into two **balanced factions**, such that:

- Links within each balanced faction are positive, and
- Links between different balanced factions are negative.

Theorem 8 (Acharya): a signed graph with adjacency matrix A is balanced if and only if A and $|A|$ have the same spectrum.

- Harary, F. (1953). "On the notion of balance of a signed graph". *Michigan Mathematical Journal*, 2(2), 143-146.
- Acharya (1980). "Spectral criterion for cycle balance in networks". *Journal of Graph Theory*, 4(1), 1-11.
- Zaslavsky (1982). "Signed graphs". *Discrete Applied Mathematics* 4, 47-74.

How close is a network to a perfectly balanced state?

How close is a network to a perfectly balanced state?

Motif-based approach: count balanced triangles or squares.

- Cartwright and Harary (1956), *Psychological Review* 63.5, 277

Frustration-based approach (spin glass theory): count frustrated edges.

- Aref and Wilson (2019), *Journal of Complex Networks* 7.2, 163–189

Dynamics-based approach: convergence to stationary state of a diffusive process as a proxy for balance.

- Kunegis et al (2010), *Proceedings of the 2010 SIAM international conference on data mining*, pp 559–570

Walk-based approach: count positive and negative walks as a proxy for cycles.

- Estrada and Benzi (2014), *Physical Review E*, 90.4, 042802
- Kirkley, Cantwell, Newman (2019), *Physical Review E* 99.1, 012320

Local (node-based) levels of balance

Node 3 has medium balance

Local (node-based) levels of balance

Node 3 has medium balance

Challenges:

- How to enumerate all cycles? → **Walk lemma**
- How to aggregate cycles of different length? → **Weight factor (inverse factorial)**

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

2) Include all walk lengths: $\sum_{k=0}^{\infty} \sum_{i=1}^N (A^k)_{ii}$

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

2) Include all walk lengths: $\sum_{k=0}^{\infty} \sum_{i=1}^N (A^k)_{ii}$

3) Penalize longer walks: $\sum_{k=0}^{\infty} \sum_{i=1}^N \frac{(A^k)_{ii}}{k!} = \text{tr}(e^A)$

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

2) Include all walk lengths: $\sum_{k=0}^{\infty} \sum_{i=1}^N (A^k)_{ii}$

3) Penalize longer walks: $\sum_{k=0}^{\infty} \sum_{i=1}^N \frac{(A^k)_{ii}}{k!} = \text{tr}(e^A)$

4) Normalize: $\kappa := \frac{\text{tr}(e^A)}{\text{tr}(e^{|A|})}$

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

2) Include all walk lengths: $\sum_{k=0}^{\infty} \sum_{i=1}^N (A^k)_{ii}$

3) Penalize longer walks: $\sum_{k=0}^{\infty} \sum_{i=1}^N \frac{(A^k)_{ii}}{k!} = \text{tr}(e^A)$

4) Normalize: $\kappa := \frac{\text{tr}(e^A)}{\text{tr}(e^{|A|})}$

Balanced graph: $\kappa = 1$
Unbalanced graph: $0 < \kappa < 1$

Balance index

1) Count balanced and unbalanced closed walks (walk lemma): $(A^k)_{ii}$

2) Include all walk lengths: $\sum_{k=0}^{\infty} \sum_{i=1}^N (A^k)_{ii}$

3) Penalize longer walks: $\sum_{k=0}^{\infty} \sum_{i=1}^N \frac{(A^k)_{ii}}{k!} = \text{tr}(e^A)$

4) Normalize:

$$\kappa := \frac{\text{tr}(e^A)}{\text{tr}(e^{|A|})}$$

Balanced graph: $\kappa = 1$

Unbalanced graph: $0 < \kappa < 1$

Also local version!

$$\kappa_i := \frac{(e^A)_{ii}}{(e^{|A|})_{ii}}$$

Diaz-Diaz, Bartesaghi, and Estrada (2024). *Journal of Applied Mathematics and Computing*, 1–24.

Up to now, we have counted **closed walks**, to quantitatively measure **balance**.

Now, we will enumerate **open walks**, to quantitatively measure **effective alliances and enmities**.

Up to now, we have counted **closed walks**, to quantitatively measure **balance**.

Now, we will enumerate **open walks**, to quantitatively measure **effective alliances and enmities**.

Nodes 3 and 4 are effective enemies (up to length 3)

Up to now, we have counted **closed walks**, to quantitatively measure **balance**.

Now, we will enumerate **open walks**, to quantitatively measure **effective alliances and enmities**.

Nodes 3 and 4 are effective enemies (up to length 3)

Communicability matrix:

$$\Gamma_{ij} = \sum_{k=0}^{\infty} \frac{(A^k)_{ij}}{k!} = (e^A)_{ij}$$

- Diaz-Diaz and Estrada. "Signed graphs in data sciences via communicability geometry." Information Sciences 710 (2025): 122096.

The distance problem in signed graphs

Definition 11: a distance is a function d that satisfies the following axioms:

- 1) Non-negativity: $d(i,j) \geq 0$
- 2) Identity of indiscernibles: $d(i,j) = 0$ iff $i = j$
- 3) Symmetry: $d(i,j) = d(j,i)$
- 4) Triangle inequality: $d(i,j) + d(j,k) \geq d(i,k)$

The distance problem in signed graphs

Definition 11: a distance is a function d that satisfies the following axioms:

- 1) Non-negativity: $d(i,j) \geq 0$
- 2) Identity of indiscernibles: $d(i,j) = 0$ iff $i=j$
- 3) Symmetry: $d(i,j) = d(j,i)$
- 4) Triangle inequality: $d(i,j) + d(j,k) \geq d(i,k)$

How can we define
a well-defined distance
on a signed graph?

Example 12: minimum-weight distance in a signed graph:

- ~~1) Non-negativity~~
- ~~2) Identity of indiscernibles~~
- ~~3) Symmetry~~
- ~~4) Triangle inequality~~

The distance problem in signed graphs

Definition 11: a distance is a function d that satisfies the following axioms:

- 1) Non-negativity: $d(i,j) \geq 0$
- 2) Identity of indiscernibles: $d(i,j) = 0$ iff $i=j$
- 3) Symmetry: $d(i,j) = d(j,i)$
- 4) Triangle inequality: $d(i,j) + d(j,k) \geq d(i,k)$

How can we define
a well-defined distance
on a signed graph?

Example 13: communicability distance in a signed graph:

$$\xi_{ij} = \sqrt{\Gamma_{ii} + \Gamma_{jj} - 2\Gamma_{ij}}$$

Very central nodes increase
the distance

Effective enemies (negative comm)
increase the distance

The distance problem in signed graphs

Definition 13: a distance is a function d that satisfies the following axioms:

- 1) Non-negativity: $d(i,j) \geq 0$
- 2) Identity of indiscernibles: $d(i,j) = 0$ iff $i=j$
- 3) Symmetry: $d(i,j) = d(j,i)$
- 4) Triangle inequality: $d(i,j) + d(j,k) \geq d(i,k)$

**ξ is a Euclidean distance,
even when the graph is
signed**

Example 13: communicability distance in a signed graph:

$$\xi_{ij} = \sqrt{\Gamma_{ii} + \Gamma_{jj} - 2\Gamma_{ij}}$$

Structural balance & signed metrics

Context X

Context Y

Mappings to unsigned networks

G

v

\mathcal{G}

v^+

v

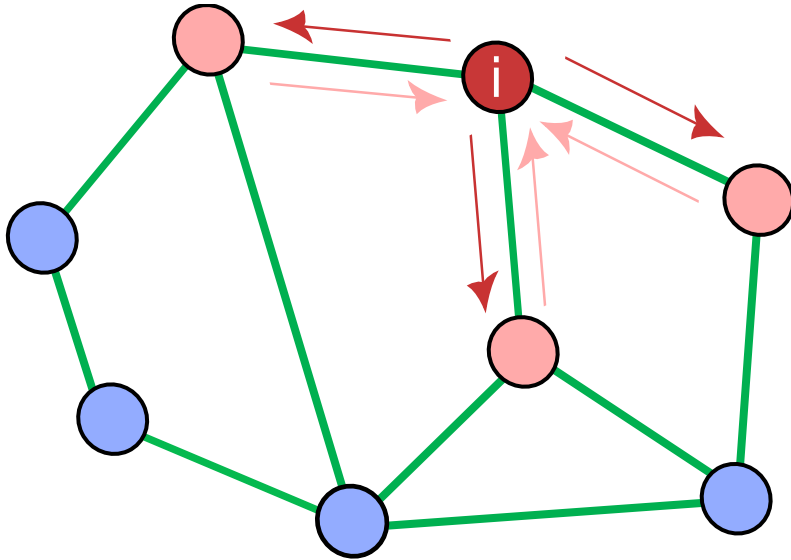
u

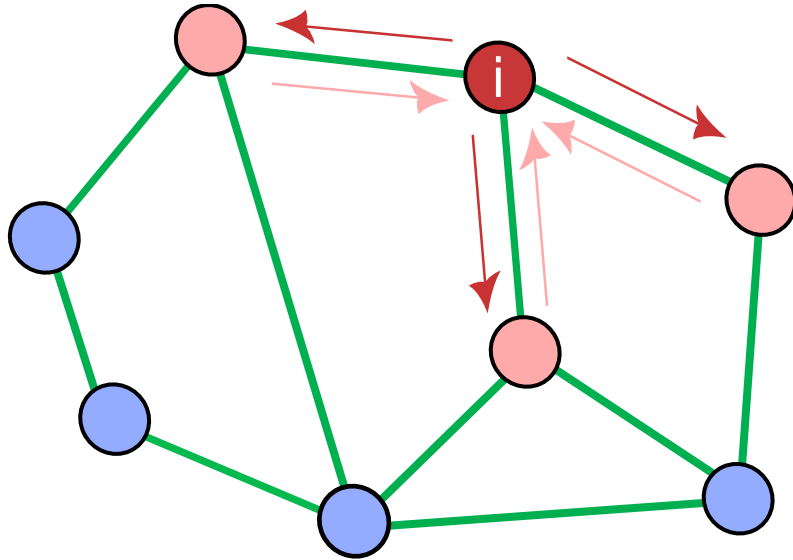
Dynamics on signed networks

Detection of communities and factions

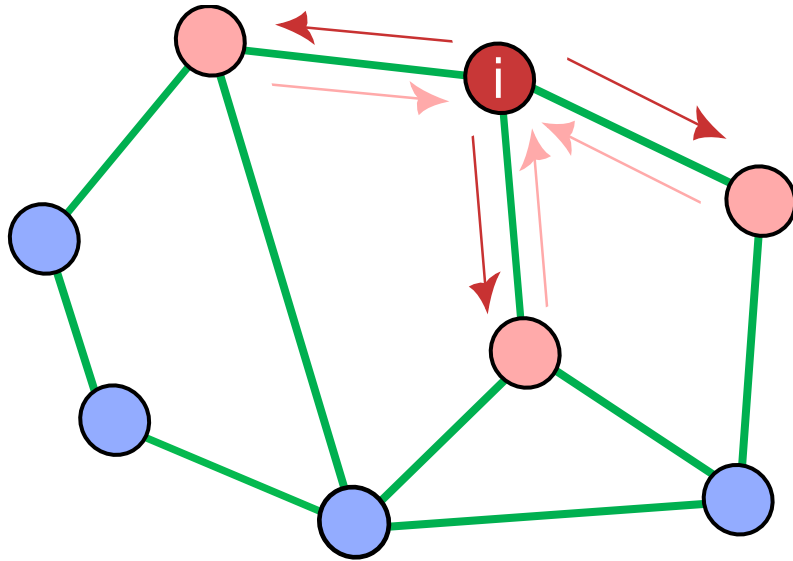
(a)

(b)



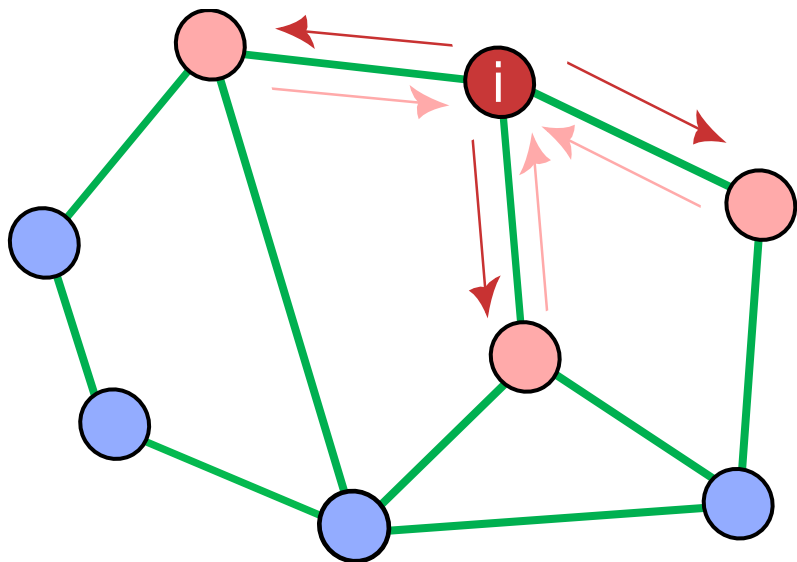


$$x_i(t+1) = x_i(t) - \sum_{j \leftarrow i} x_i(t) + \sum_{j \rightarrow i} x_j(t)$$



$$x_i(t+1) = x_i(t) - \sum_{j \leftarrow i} x_i(t) + \sum_{j \rightarrow i} x_j(t)$$

$$\dot{x}(t) = -(K - A)x(t)$$



$$x_i(t+1) = x_i(t) - \sum_{j \leftarrow i} x_i(t) + \sum_{j \rightarrow i} x_j(t)$$

$$\dot{x}(t) = -(K - A)x(t)$$

$$\dot{x}(t) = -Lx(t)$$

**Diffusion equation
for graphs**

Definition 14: the **Laplacian operator** L of an unsigned network is given by:

$$L_{ij} = \begin{cases} k_i & \text{if } i = j \\ -A_{ij} & \text{otherwise} \end{cases} \quad \text{where } k_i = \sum_j A_{ij}$$

Definition 15: the **Laplacian operator** L of a signed network is given by:

$$L_{ij} = \begin{cases} k_i & \text{if } i = j \\ -A_{ij} & \text{otherwise} \end{cases} \quad \text{where } k_i = \sum_j |A_{ij}|$$

Definition 15: the **Laplacian operator** L of a signed network is given by:

$$L_{ij} = \begin{cases} k_i & \text{if } i = j \\ -A_{ij} & \text{otherwise} \end{cases} \quad \text{where } k_i = \sum_j |A_{ij}|$$

Theorem 16: the Laplacian operator of a signed graph is a **positive semidefinite**. Moreover, L has a null eigenvalue if and only if the graph is balanced.

Linear dynamics (**diffusion**):

$$\dot{x}(t) = -Lx(t) \quad \Longrightarrow \quad x(t) = e^{-Lt}x_0$$

Linear dynamics (**diffusion**):

$$\dot{x}(t) = -Lx(t) \quad \Longrightarrow \quad x(t) = e^{-Lt}x_0$$

Eigenmode expansion:

$$x_i(t) = \sum_k e^{-\lambda_k t} \psi_k(i) \psi_k(j) x_0(j)$$

Linear dynamics (**diffusion**):

$$\dot{x}(t) = -Lx(t) \quad \Longrightarrow \quad x(t) = e^{-Lt}x_0$$

Eigenmode expansion:

$$x_i(t) = \sum_k e^{-\lambda_k t} \psi_k(i) \psi_k(j) x_0(j)$$

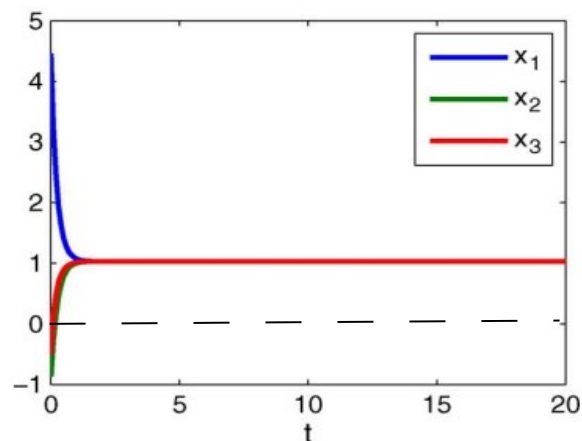
In the stationary state:

$$\lim_{t \rightarrow \infty} x_i(t) = \begin{cases} 0 & \text{if } \lambda_0 > 0 \\ C\psi_0 & \text{otherwise} \end{cases} \quad (\psi_0)_i = \begin{cases} +1 & \text{if } i \in G_1 \\ -1 & \text{if } i \in G_2 \end{cases}$$

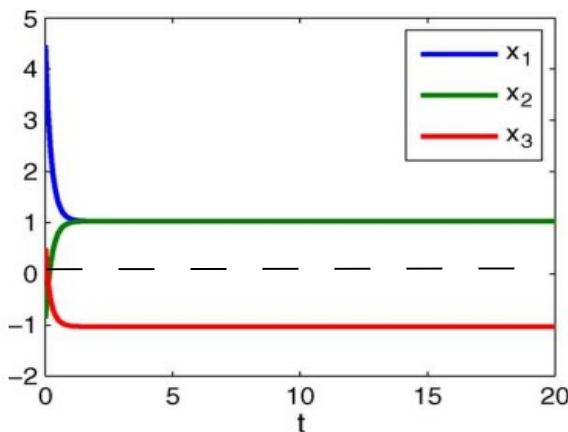
Theorem 17: the stationary state of a diffusive process on a signed network depends on the structural balance of the network. In particular, the stationary state is:

- **Consensus** if the network is unsigned.
- **Agreed dissensus** if the network is balanced.
- **Absence of opinions** if the network is unbalanced.

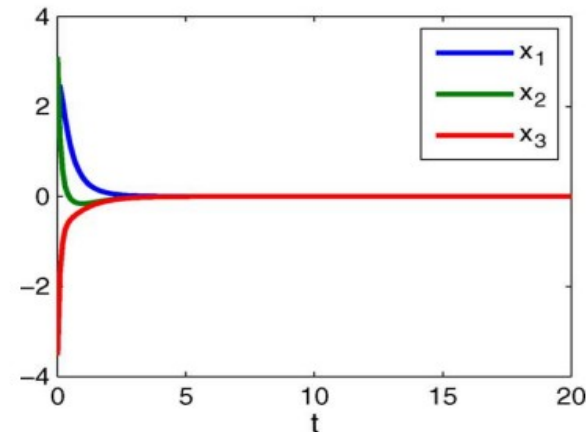
Consensus



Agereed dissensus

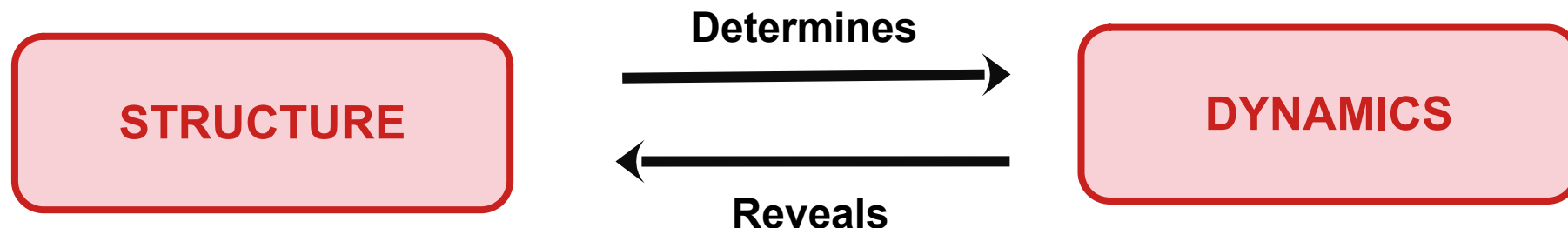


Absence of opinions



Theorem 17: the stationary state of a diffusive process on a signed network depends on the structural balance of the network. In particular, the stationary state is:

- **Consensus** if the network is unsigned.
- **Agreed dissensus** if the network is balanced.
- **Absence of opinions** if the network is unbalanced.



Mathematics of signed graphs:

- Thomas Zaslavsky. “Signed Graphs”. Discrete Applied Mathematics 4 (1982), pages 47–74.
- Thomas Zaslavsky. Matrices in the Theory of Signed Simple Graphs. 2013. arXiv: 1303.3083 [math].

Structural balance:

- Frank Harary. “On the Notion of Balance of a Signed Graph”. Michigan Mathematical Journal 2 (1953), pages 143–146
- Dorwin Cartwright and Frank Harary. “Structural Balance: A Generalization of Heider’s Theory.” Psychological Review 63.5 (1956), page 277.

Balance indices:

- Ernesto Estrada and Michele Benzi. “Walk-Based Measure of Balance in Signed Networks: Detecting Lack of Balance in Social Networks”. Physical Review E 90.4 (2014), page 042802.

Laplacian and dynamics:

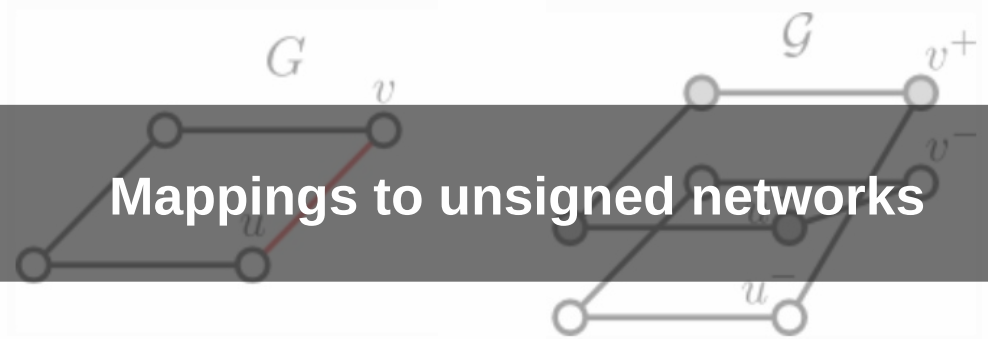
- Jérôme Kunegis, Stephan Schmidt, Andreas Lommatzsch, Jürgen Lerner, Ernesto W De Luca, and Sahin Albayrak. “Spectral analysis of signed graphs for clustering, prediction and visualization”. Proceedings of the 2010 SIAM international conference on data mining. SIAM. 2010, pages 559–570.
- Claudio Altafini. “Consensus problems on networks with antagonistic interactions”. IEEE transactions on automatic control 58.4 (2012), pages 935–946.

Structural balance & signed metrics

Context X

Context Y

Mappings to unsigned networks

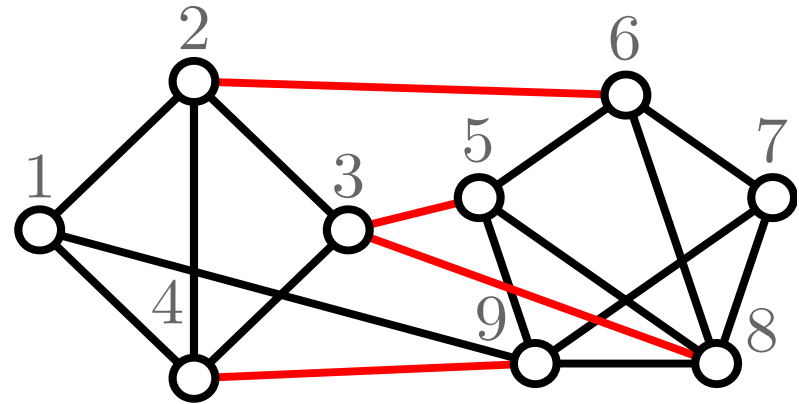
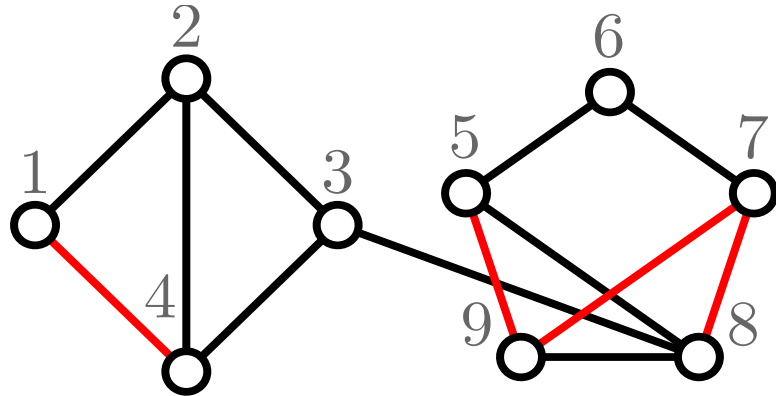


Dynamics on signed networks

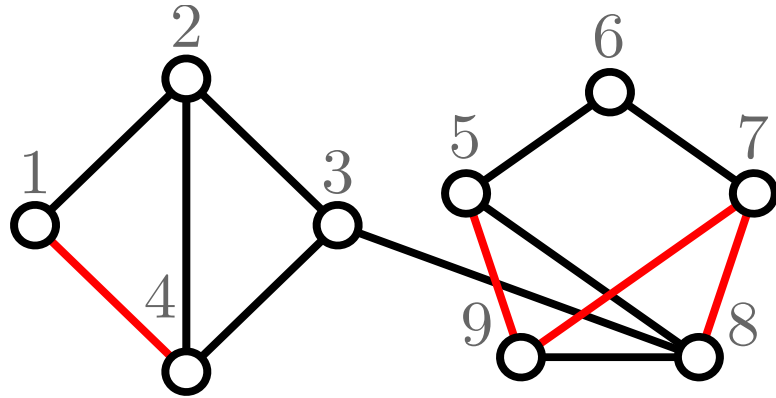
Detection of communities and factions



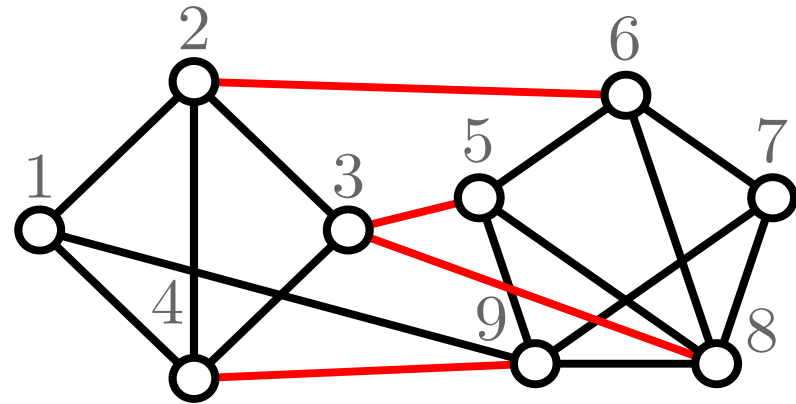
Q: which signed network has mesoscale structure?



Q: which signed network has mesoscale structure?

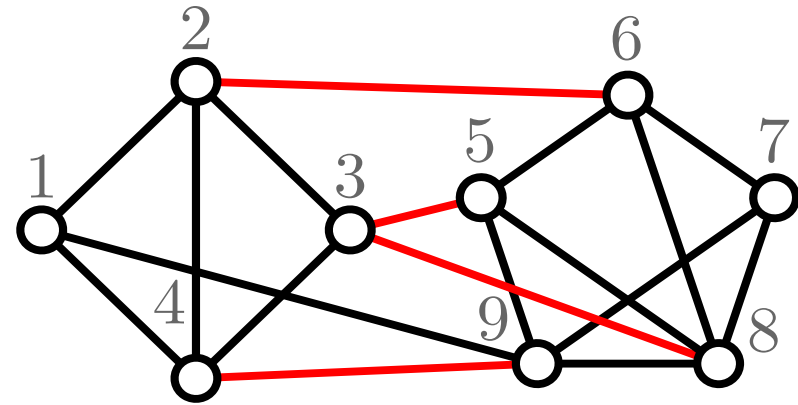
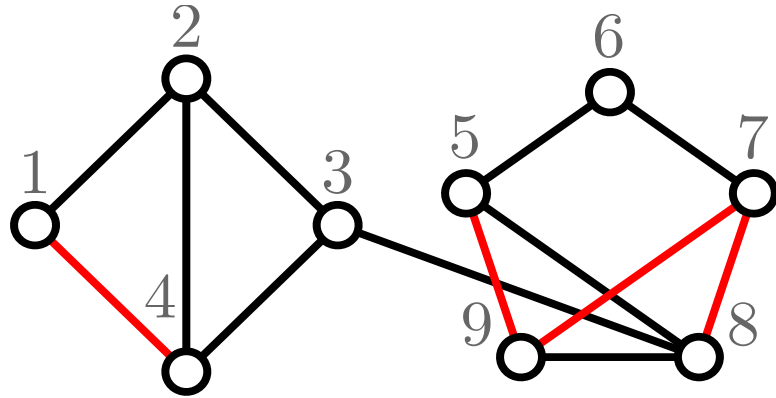


COMMUNITIES



FACTIONS

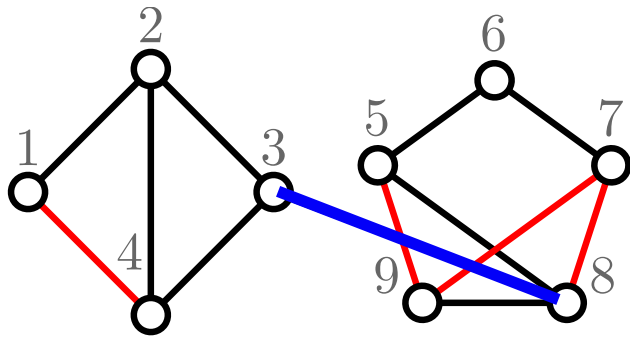
Q: which signed network has mesoscale structure?



Definition (Cut set)

For a partition $V = U_1 \cup U_2$, the cut-set $C(U_1, U_2)$ is the set of edges with one endpoint in U_1 and the other in U_2 .

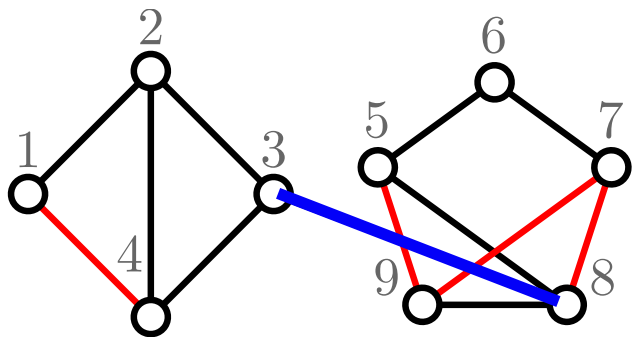
The edge connectivity $\kappa_e(G)$ is the size of a smallest cut-set of G .



Communities have small edge connectivity

Definition (Cut set)

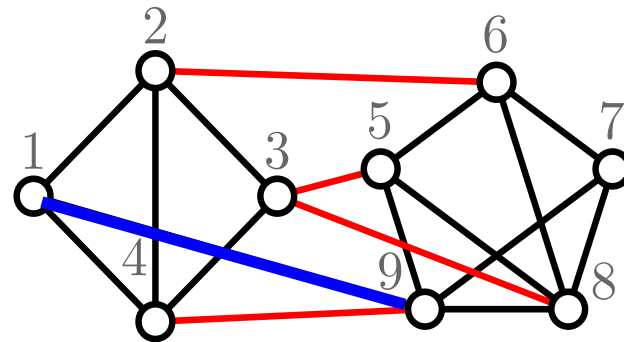
For a partition $V = U_1 \cup U_2$, the cut-set $C(U_1, U_2)$ is the set of edges with one endpoint in U_1 and the other in U_2 . The edge connectivity $\kappa_e(G)$ is the size of a smallest cut-set of G .



Communities have small edge connectivity

Definition (Frustration set)

A frustration set is a set of edges whose removal makes the signed graph balanced. The frustration index $\phi(G)$ is the size of a smallest frustration set of G .

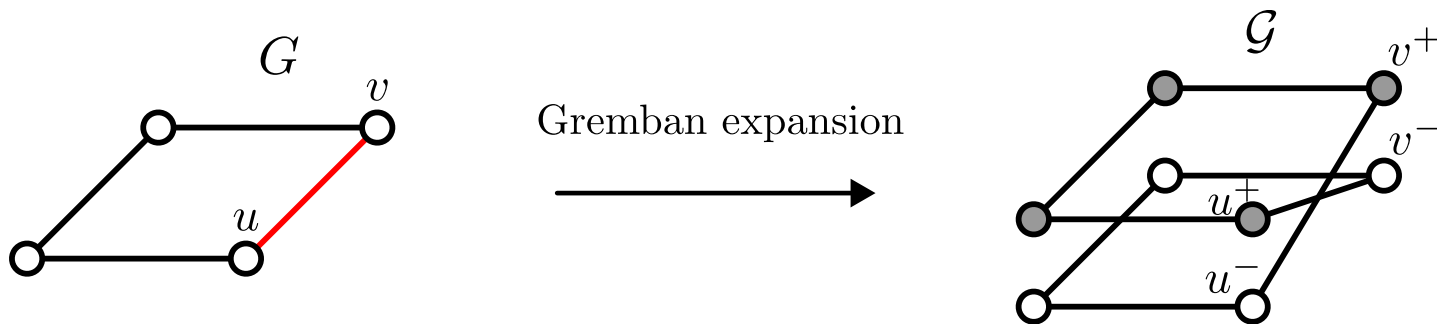


Factions have small frustration index

Definition (Gremban expansion)

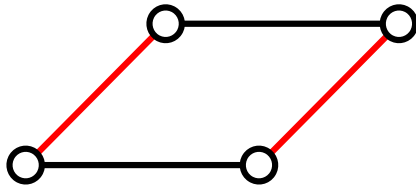
The *Gremban expansion* of a signed graph G is the unsigned graph \mathcal{G} with $2n$ vertices and $2m$ edges, defined as follows:

- Each node v gets mapped to two polarities v^+ and v^- .
- Positive links (u, v) get mapped to (u^+, v^+) and (u^-, v^-) .
- Negative links (u, v) get mapped to (u^+, v^-) and (u^-, v^+) .

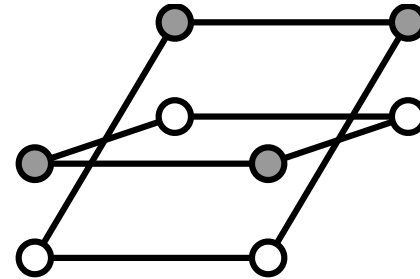


Theorem (1)

A connected signed graph is balanced iff its Gremban expansion is disconnected.

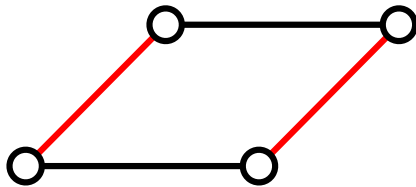


Gremban expansion

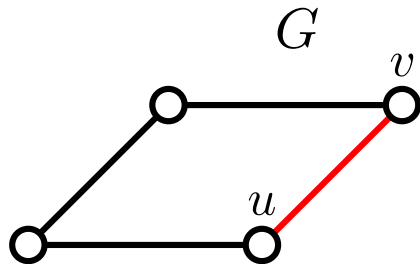
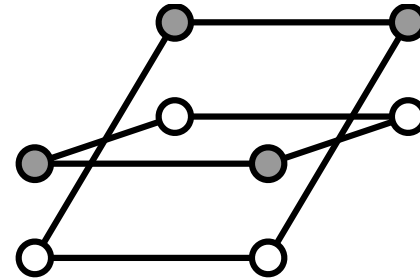


Theorem (1)

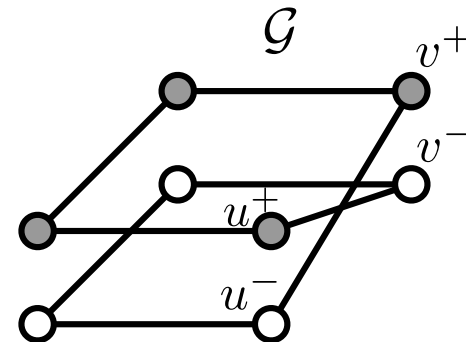
A connected signed graph is balanced iff its Gremban expansion is disconnected.



Gremban expansion



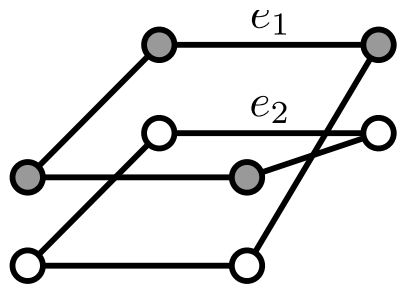
Gremban expansion



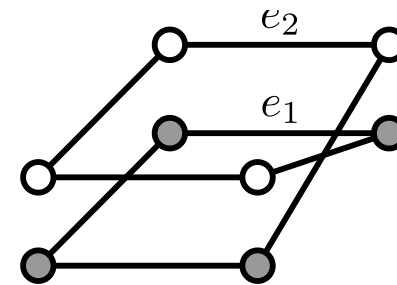
Definition (Gremban involution)

The *Gremban involution* η swaps the two polarities of every node:

$$\eta(v^+) = v^-, \quad \eta(v^-) = v^+.$$



Gremban involution η



Definition (Gremban involution)

The *Gremban involution* η swaps the two polarities of every node:

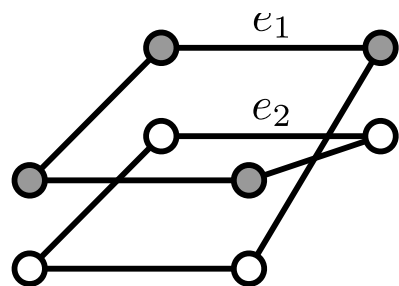
$$\eta(v^+) = v^-, \quad \eta(v^-) = v^+.$$

Definition (Gremban symmetry)

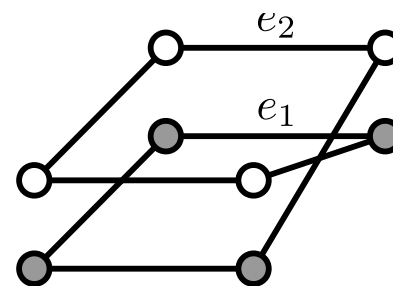
- A set (of nodes, edges, or a subgraph) is *Gremban-symmetric* if it is invariant under the involution η , i.e.

$$\eta(\mathcal{X}) = \mathcal{X}.$$

- A bipartition $(\mathcal{U}_1, \mathcal{U}_2)$ of $V(\mathcal{G})$ is *Gremban-symmetric* if $\eta(\mathcal{U}_1) = \mathcal{U}_1$ or $\eta(\mathcal{U}_1) = \mathcal{U}_2$.



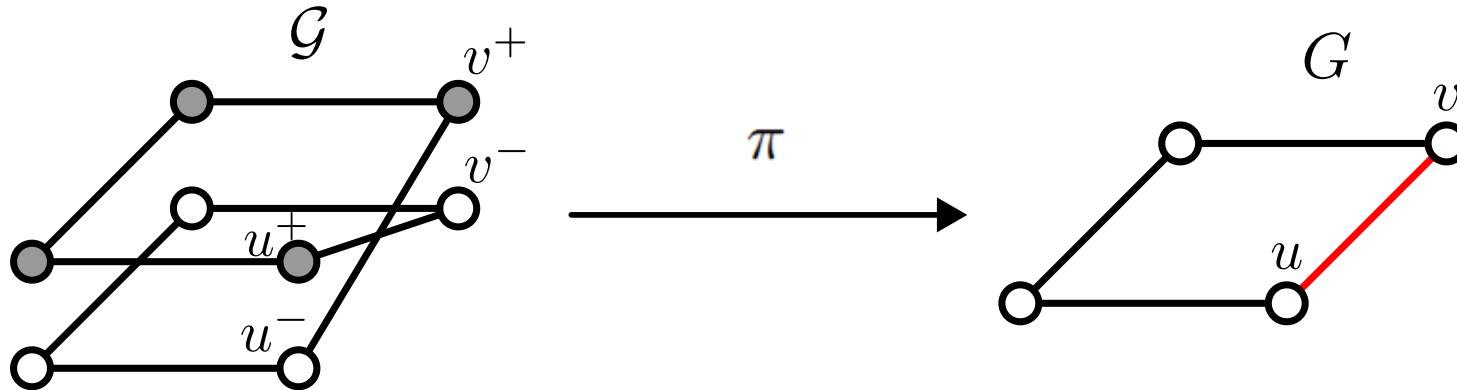
Gremban involution η



Definition (Projection map)

The *projection* π sends each polarity back to its node:

$$\pi(v^+) = \pi(v^-) = v.$$



So far, we have:

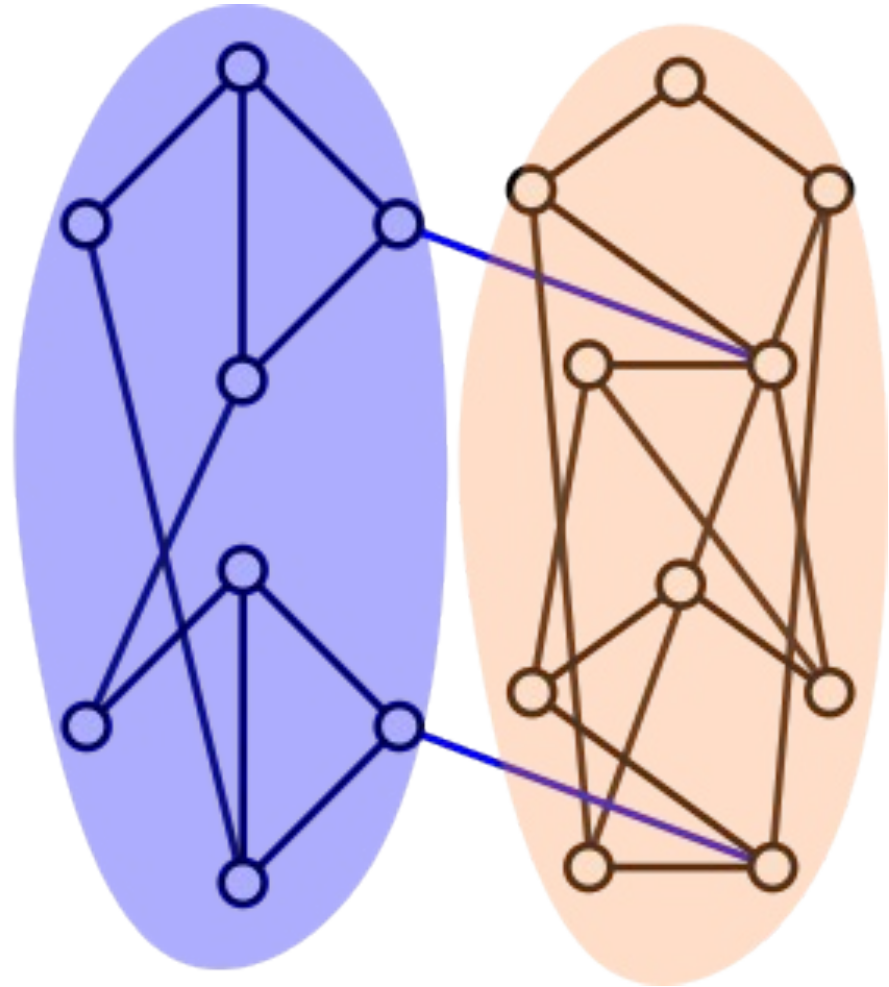
- ① A way to operationalize communities and factions (small **cut-sets** and small **frustration sets**).
- ② An operation, the **Gremban expansion**, that maps a signed graph to an unsigned one.
- ③ A notion of **symmetry** in the expanded space (invariance under the **involution** η).
- ④ A way of **projecting** back to the original space, but only for Gremban-symmetric objects.

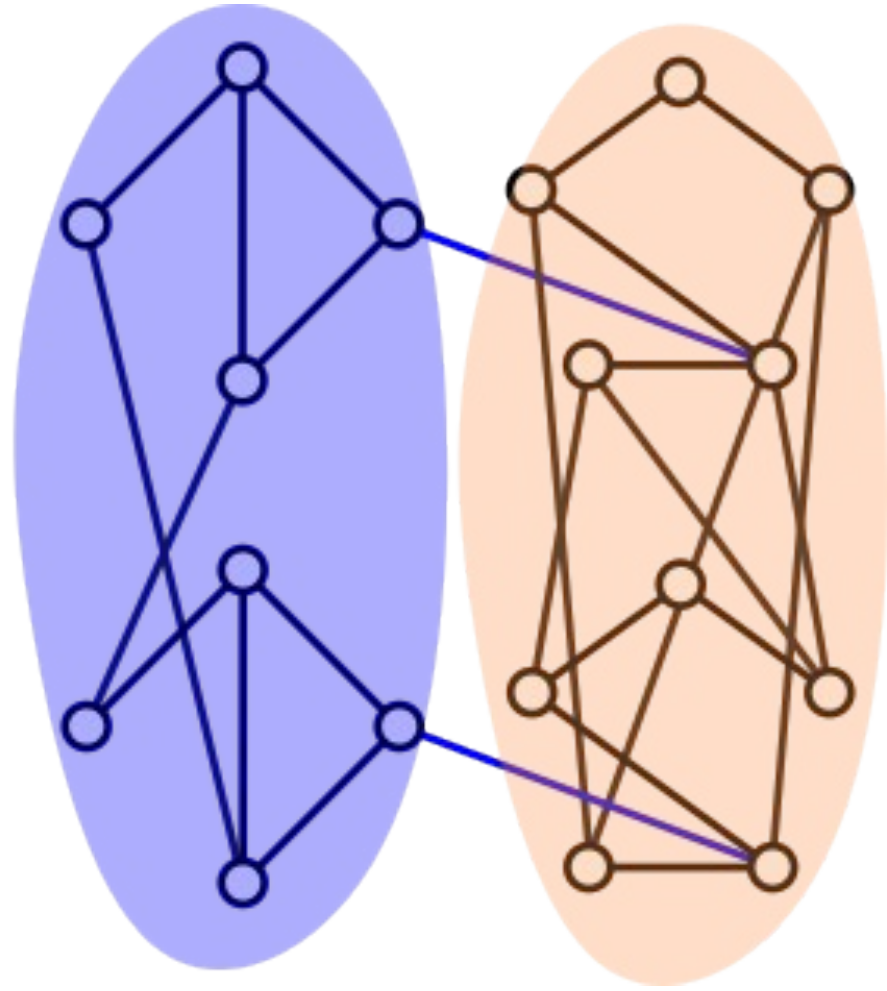
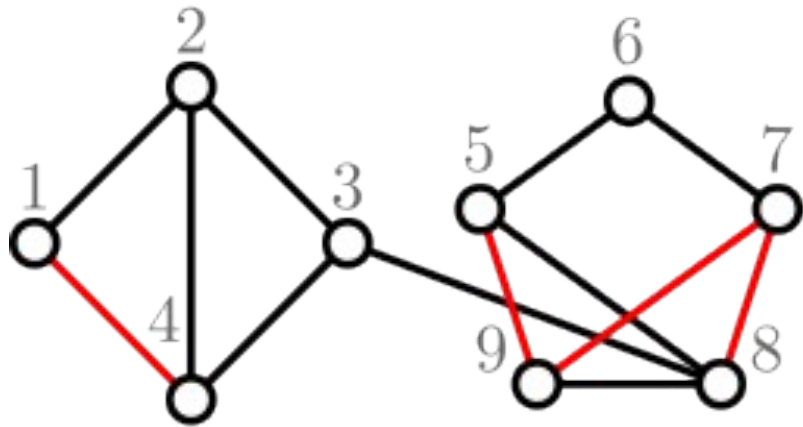
What types of structures structure in the expanded space correspond to communities and factions?

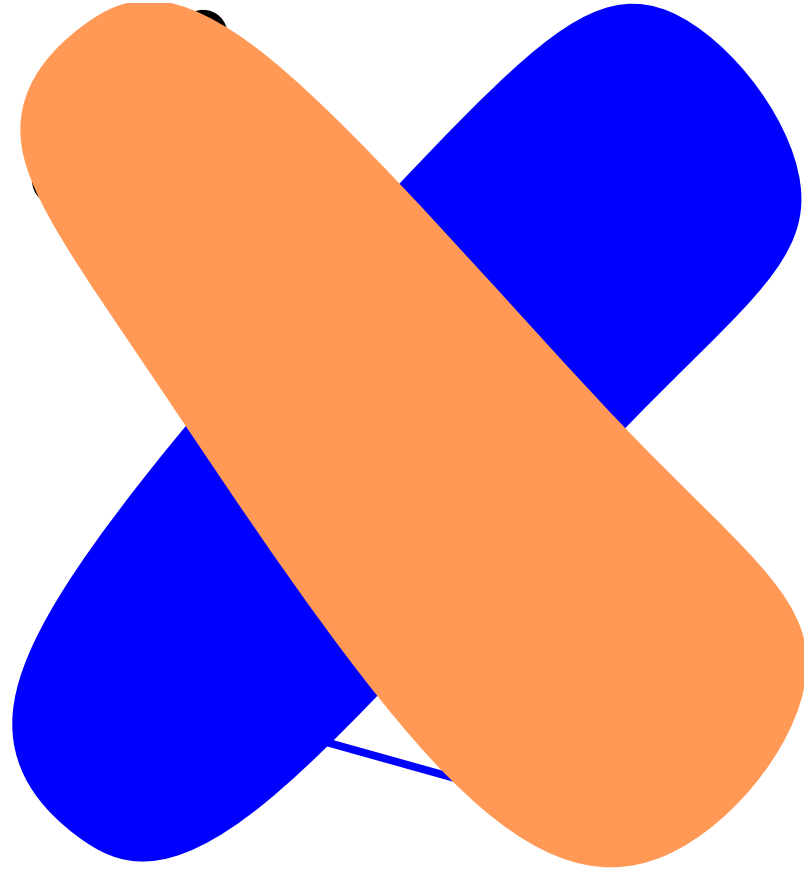
Theorem (3)

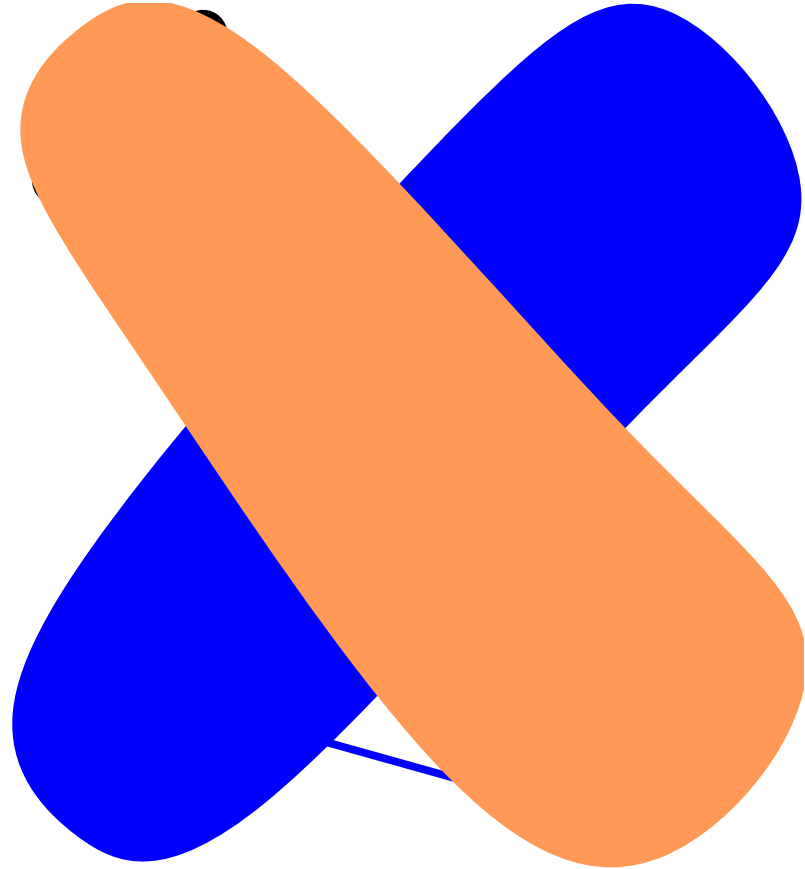
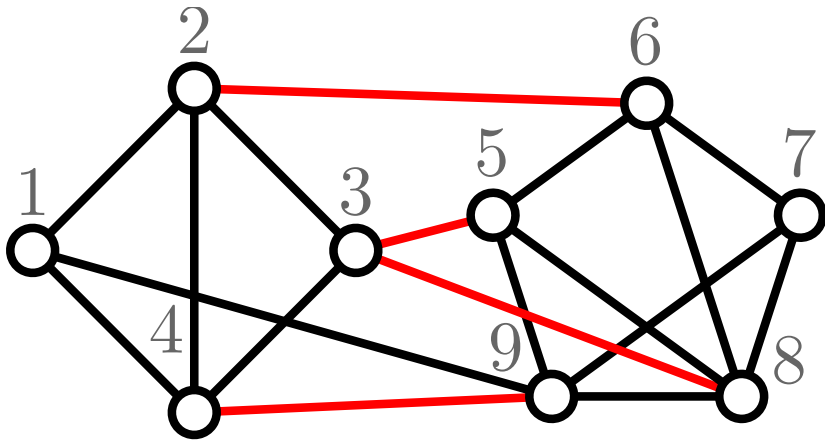
Every cut or frustration set in G lifts to a symmetric cut in \mathcal{G} . Conversely, every Gremban-symmetric bipartition $(\mathcal{U}_1, \mathcal{U}_2)$ of \mathcal{G} projects to:

$$\begin{cases} \text{cut-sets in } G & \text{if } \eta(\mathcal{U}_1) = \mathcal{U}_1, \\ \text{frustration sets in } G & \text{if } \eta(\mathcal{U}_1) = \mathcal{U}_2. \end{cases}$$



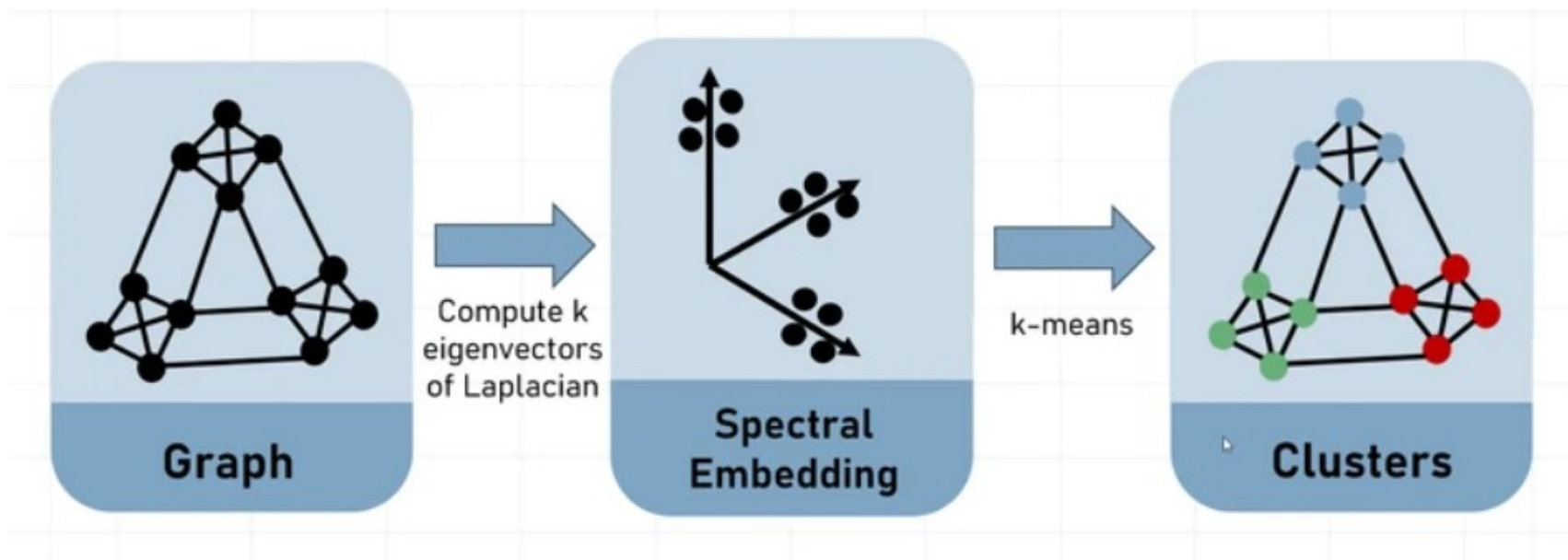






Spectral Clustering Principle:

Low eigenvectors of the Laplacian have similar values on well-connected nodes → **eigenvectors of L reveal community structure**



- Ulrike von Luxburg. A tutorial on spectral clustering. Statistics and Computing, 17(4):395–416, 2007.
- Image: Macgregor, Peter. "Fast and simple spectral clustering in theory and practice." Advances in Neural Information Processing Systems 36 (2023): 34410-34425.

① Signed adjacency matrix.

$$A = A^+ - A^-$$

① Signed adjacency matrix.

$$A = A^+ - A^-$$



② Gremban-expanded adjacency matrix (**non-negative!**).

$$\mathcal{A} = \begin{pmatrix} A^+ & A^- \\ A^- & A^+ \end{pmatrix}$$

① Signed adjacency matrix.

$$A = A^+ - A^-$$



② Gremban-expanded adjacency matrix (**non-negative!**).

$$\mathcal{A} = \begin{pmatrix} A^+ & A^- \\ A^- & A^+ \end{pmatrix}$$



$$L = K - A$$

③ Laplacian of the expanded adjacency matrix.

$$\mathcal{L} = \begin{pmatrix} K - A^+ & -A^- \\ -A^- & K - A^+ \end{pmatrix}$$

Theorem (4)

Let $L = K - A$ be the signed Laplacian and $\bar{L} = K - |A|$ the unsigned Laplacian. Then, $\mathcal{L} \sim \bar{L} \oplus L$.

$$\mathcal{L} = \mathcal{U} \begin{pmatrix} \bar{L} & 0 \\ 0 & L \end{pmatrix} \mathcal{U}^\top$$

- $(\lambda, x) \in L \Rightarrow (\lambda, (x, -x)) \in \mathcal{L},$
- $(\mu, y) \in \bar{L} \Rightarrow (\mu, (y, y)) \in \mathcal{L}.$

Theorem (4)

Let $L = K - A$ be the signed Laplacian and $\bar{L} = K - |A|$ the unsigned Laplacian. Then, $\mathcal{L} \sim \bar{L} \oplus L$.

$$\mathcal{L} = \mathcal{U} \begin{pmatrix} \bar{L} & 0 \\ 0 & L \end{pmatrix} \mathcal{U}^\top$$

- $(\lambda, x) \in L \Rightarrow (\lambda, (x, -x)) \in \mathcal{L},$
- $(\mu, y) \in \bar{L} \Rightarrow (\mu, (y, y)) \in \mathcal{L}.$

Communities → Unsigned topology → Unsigned Laplacian → Symmetric eigenvectors
→ Gremban-symmetric node partitions

Factions → Signed topology → Signed Laplacian → Antisymmetric eigenvectors
→ Gremban-antisymmetric node partitions

Algorithm:

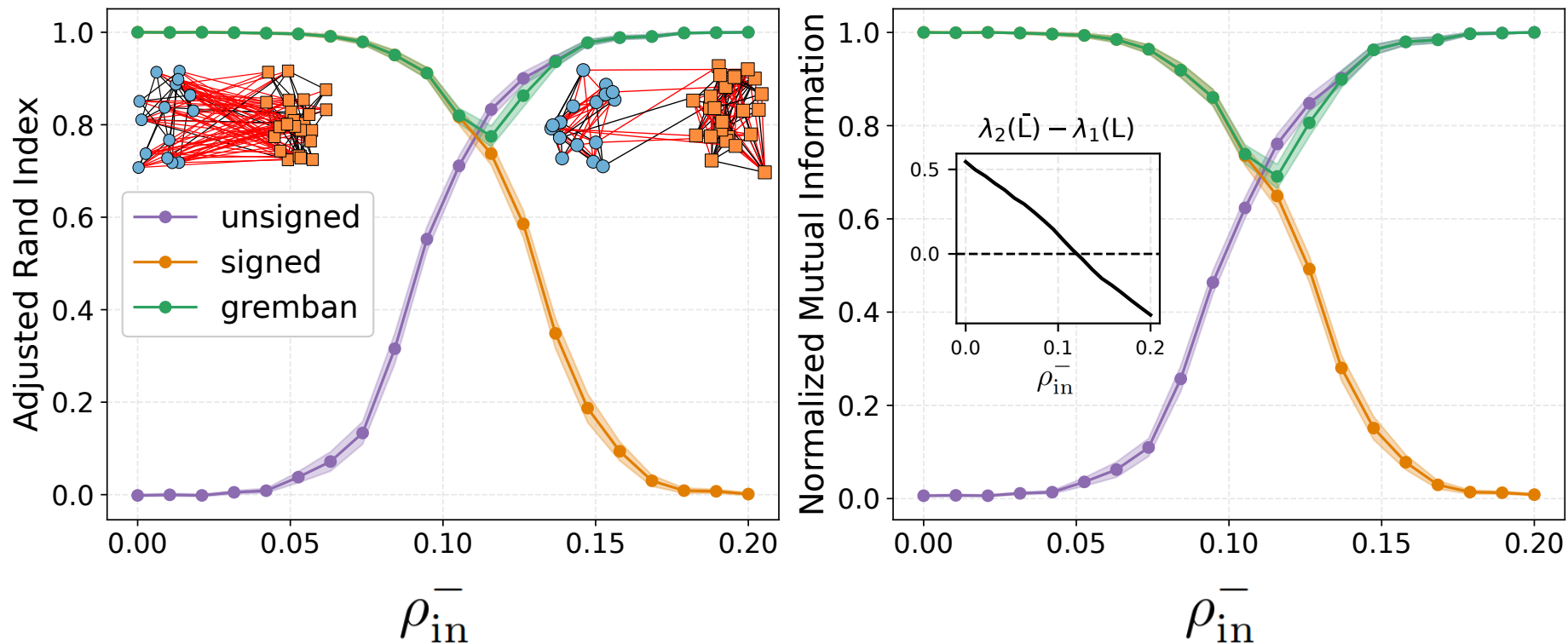
- ① Expand signed graph $G \mapsto \mathcal{G}$.
- ② Compute Laplacian \mathcal{L} of \mathcal{G} .
- ③ Extract first non-constant $k - 1$ eigenvectors $\{\psi_2, \dots, \psi_k\}$.
- ④ Embed nodes in \mathbb{R}^{k-1} and run k -means.
- ⑤ Interpret clusters:
 - Symmetric \rightarrow communities
 - Antisymmetric \rightarrow factions inside communities

Algorithm:

- ① Expand signed graph $G \mapsto \mathcal{G}$.
- ② Compute Laplacian \mathcal{L} of \mathcal{G} .
- ③ Extract first non-constant $k - 1$ eigenvectors $\{\psi_2, \dots, \psi_k\}$.
- ④ Embed nodes in \mathbb{R}^{k-1} and run k -means.
- ⑤ Interpret clusters:
 - Symmetric \rightarrow communities
 - Antisymmetric \rightarrow factions inside communities

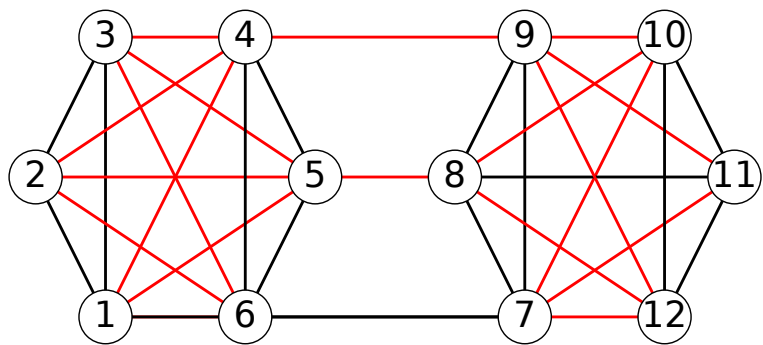
This detects **both communities and factions** and disentangles them in a principled way!

Spectral clustering in the Gremban expansion

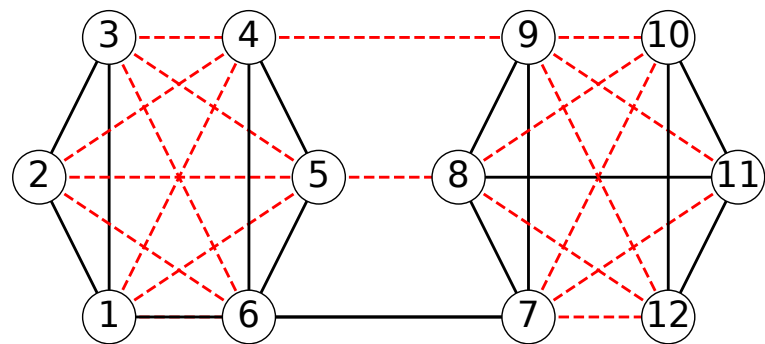


- Diaz-Diaz, Fernando, Karel Devriendt, and Renaud Lambiotte. "Gremban Expansion for Signed Networks: Algebraic and Combinatorial Foundations for Community-Faction Detection." arXiv preprint arXiv:2509.14193 (2025).
- Fox, Manteuffel, and Sanders. Numerical methods for Gremban's expansion of signed graphs. SIAM Journal on Scientific Computing, 39(5):S945–S968, 2017

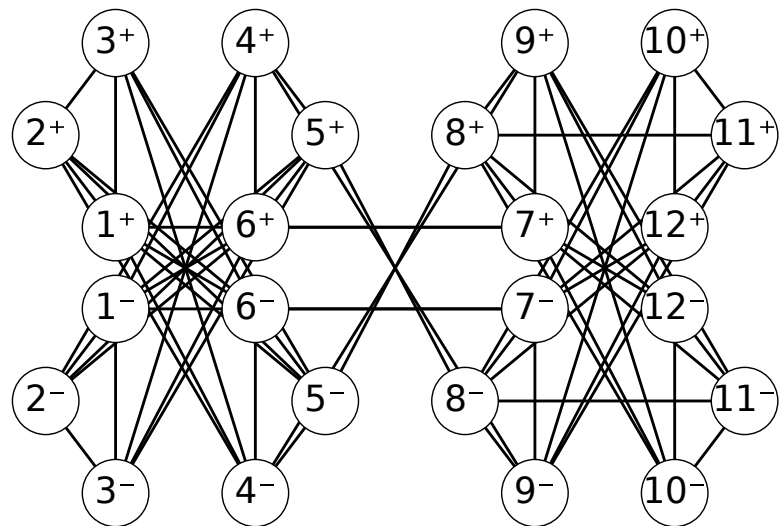
(a)



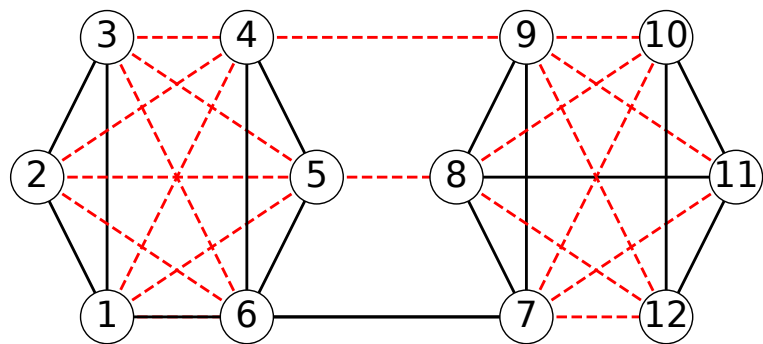
(a)



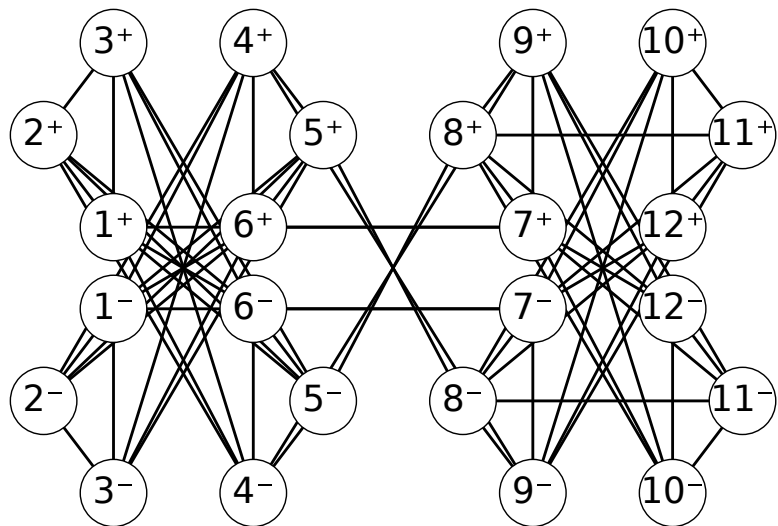
(b)



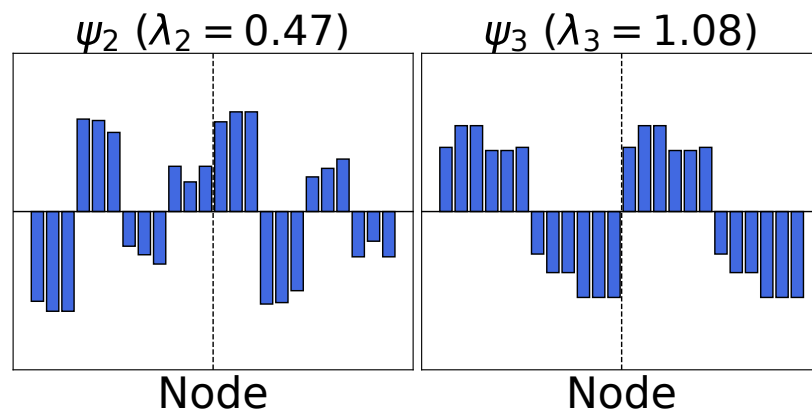
(a)



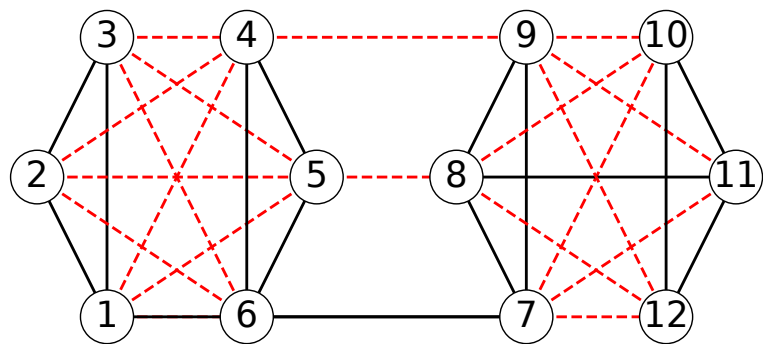
(b)



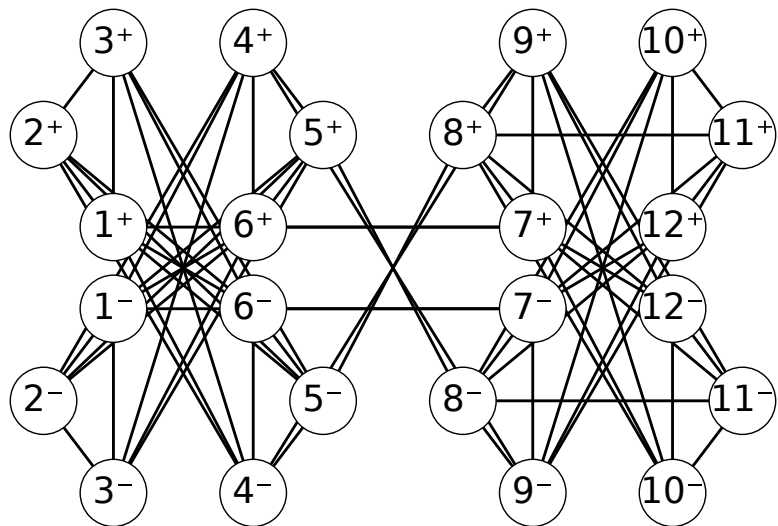
(c)



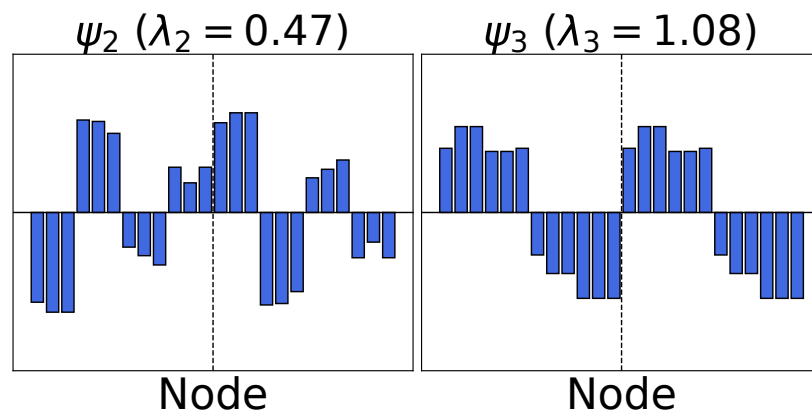
(a)



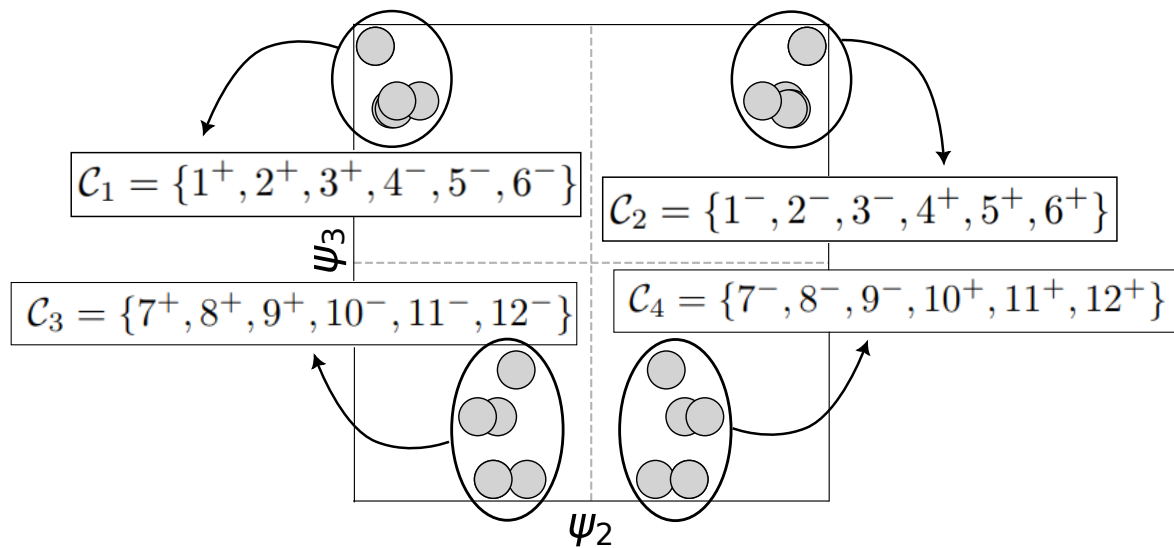
(b)



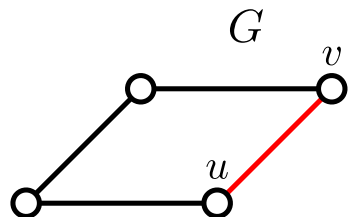
(c)



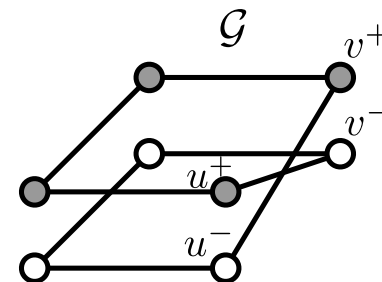
(d)



- Signed graphs can show two mesoscale structures: **communities and factions**.
- Using the **Gremban expansion**, we map a signed graph to an unsigned one without losing sign information.
- Communities in the expanded graph map back to **communities or factions** in the original, depending on the **symmetry** of the partition.



Gremban expansion



Results available in:

arXiv

> cs > arXiv:2509.14193

Search

Computer Science > Discrete Mathematics

[Submitted on 17 Sep 2025]

**Gremban Expansion for Signed Networks:
Algebraic and Combinatorial Foundations for
Community-Faction Detection**

Fernando Diaz-Diaz, Karel Devriendt, Renaud Lambiotte

Thank you!

Thank you for your attention

Contact: fddiaz@math.uc3m.es

Gremban expansion:



Signet review:



Signed Networks: theory, methods, and applications

Fernando Diaz-Diaz^{1,2}, Elena Candellone^{3,4}, Miguel A. González-Casado¹,
Emma Fraxanet⁵, Antoine Vendeville^{6,7,8}, Irene Ferri^{9,10,11}, and Andreia
Sofia Teixeira^{12,13}

¹Universidad Carlos III de Madrid, Departamento de Matemáticas, Grupo Interdisciplinar de Sistemas
Complejos, 28911 Leganés, Spain

²Institute of Cross-Disciplinary Physics and Complex Systems, IFISC (UIB-CSIC), 07122 Palma de Mallorca,
Spain

³Department of Methodology and Statistics, Utrecht University, Utrecht, Netherlands

⁴Centre for Complex Systems Studies, Utrecht University, Utrecht, Netherlands

⁵Department of Engineering, Universitat Pompeu Fabra, Barcelona 08018, Spain

⁶Sciences Po médialab, 75007 Paris, France

⁷Complex Systems Institute of Paris Ile-de-France CNRS, 75013 Paris, France

⁸Learning Planet Institute, Learning Transitions unit, CY Cergy Paris University, 75004 Paris, France

⁹Departament de Física de la Matèria Condensada, Universitat de Barcelona (UB), c. Martí i Franquès, 1, 08028
Barcelona, Spain

¹⁰Institut de Recerca en Sistemes Complexos (UBICS), Universitat de Barcelona (UB), Barcelona, Spain

¹¹The Roux Institute, Network Science Institute, Northeastern University, 04101 Portland, ME USA

¹²BRAN Lab, Network Science Institute, Northeastern University London, London E1W 1LP, United Kingdom

¹³LASIGE, Departamento de Informática, Universidade de Lisboa, Campo Grande 1749-016, Lisboa, Portugal